Data processing device, method of operating a data processing device and method for
compiling a program

The present invention relates to a data processing device.

The invention further relates to a method of operating a data processing
device.

The invention further relates to a method for compiling a program.

5

Modern signal processing systems are designed to support multiple standards
and to provide high performance. Multimedia and telecom are typical areas where such
combined requirements can be found. The need for high performance leads to architectures
10   that may include application specific hardware accelerators. In the HW/SW co-design
community, "mapping" refers to the problem of assigning the functions of the application
program to a set of operations that can be executed by the available hardware components
[1][2]. Operations may be arranged in two groups according to their complexity: fine-grain
and coarse-grain operations.

15   Examples of fine-grain operations are addition, multiplication, and conditional
jump. They are performed in a few clock cycles and only a few input values are processed at
a time. Coarse-grain operations process a bigger amount of data and implement a more
complex functionality such as FFT-butterfly, DCT, or complex multiplication.

A hardware component implementing a coarse-grain operation is characterized
20   by a latency that ranges from few cycles to several hundreds of cycles. Moreover, data
consumed and produced by the unit is not concentrated at the end and at the beginning of the
course grain operation. On the contrary, data communications to and from the unit are
distributed during the execution of the whole course grain operation. Consequently, the
functional unit exhibits a (complex) *timeshape* in terms of Input-Output behavior [9].
25   According to the granularity (coarseness) of the operations, architectures may be grouped in
two different categories, namely processor architectures and heterogeneous multi-processor
architectures, defined as follows:

-        Processor architectures: The architecture consists of a heterogeneous
collection of Functional Units (FUs) such as ALUs and multipliers. Typical architectures in

this context are general-purpose CPU and DSP architectures. Some of these, such as VLIW and superscalar architectures can have multiple operations executed in parallel. The FUs execute fine-grain operations and the data has typically a "word" grain size.

-        Heterogeneous multi-processor architectures: The architecture is made of

5   dedicated Application Specific Instruction set Processors (ASIPs), ASICs and standard DSPs and CPUs, connected via busses. The hardware executes coarse-grain operations such as a 256 input FFT, hence data has a "block of words" grain size. In this context, operations are often regarded as tasks or processes.

        The two architectural approaches above described are always been kept

10   separated.

        It is a purpose of the invention to provide a data processing device wherein a (co)-processors are embedded as FUs in a VLIW processor datapath, wherein the VLIW

15   processor can have FUs executing operations having different latencies and working on a variety of data granularities at the same time.

        It is a further purpose of the invention to provide a method for operating such a data processing device.

        It s a further purpose of the invention to provide a method for compiling a

20   program which efficiently schedules a mixture of fine-grain and coarse-grain operations, minimizing schedule's length and VLIW instruction width.

        A data processing device according to the invention at least comprises a master controller, a first functional unit which includes a slave controller, a second functional unit, which functional units share common memory means, the device being programmed for

25   executing an instruction by the first functional unit, the execution of said instruction involving input/output operations by the first functional unit, wherein output data of the first functional unit is processed by the second functional unit during said execution and/or the input data is generated by the second functional unit during said execution.

        The first functional unit is for example Application Specific Instruction set

30   Processor (ASIP), an ASIC, a standard DSP or a CPU. The second functional unit typically executes fine-grain operations such as an ALU or a multiplier. The common memory means shared by the first and the second unit may be a program memory which comprises the instructions to be carried out by these units. Otherwise the common memory means may be used for data storage.

Introducing coarse-grain operations has a beneficial influence on the microcode width. Firstly, because FUs executing coarse-grain operations have internally their own controller. Therefore, the VLIW controller needs less instruction bits to steer the entire datapath. Secondly, exploiting the I/O timeshape makes it possible to deliver and consume data even if the operation itself is not completed, hence shortening signals' lifetime and, therefore, the number of datapath registers. The instruction bits needed to address datapath registers and steering in parallel a large number of datapath resources are two important factors contributing to the large width of the VLIW microcode. Ultimately, enhancing the instruction level parallellism (ILP) has a positive influence on the schedule length, and hence, on microcode length. Keeping microcode area small is an essential requisite for embedded applications aiming at high performances and coping with long and complex program codes. The internal schedule of the FUs will be partially taken into account while scheduling the application. In this way, a FU's *internal* schedule could be considered as *embedded* in the application's VLIW schedule. Doing so, the knowledge on the I/O timeshape might be exploited to provide or withdraw data from the FU in a "just in time" fashion. The operation can start even if not all data consumed by the unit is available. A FU performing coarse-grain operations can be re-used as well. This means that it can be maintained in the VLIW datapath, while the actual *use* of its output data will be different.

It is remarked that commercially available DSPs, based on the VLIW architecture are known which limit the complexity of custom operations executed by the datapath's FUs. The R.E.A.L. DSP [3], for instance, allows the introduction of custom units, called Application-specific eXecution Units (AXU). However, the latency of these functional units is limited to one clock cycle. Other DSPs like the TI 'C6000 [4] may contain FUs with latency ranging from one to four cycles. The Philips Trimedia VLIW architecture [5] allows multi-cycle and pipelined operation ranging from one to three cycles. The architectural level synthesis tool Phideo [10] can handle operations with timeshapes, but is not suited for control-dominated applications. Mistral2 [11] allows the definition of timeshape under the restriction that signals are passed to separate I/O ports of the FU. Currently, no scheduler can cope well with FUs with complex timeshapes. To simplify the scheduler's job, the unit performing a coarse-grain operation is traditionally characterized only by its latency and the operation is regarded as atomic. Consequently, this approach lengthens the schedule because all data must be available before starting the operation, regardless the fact that the unit could already perform some of its computations without having the total amount of input data. This approach lengthens the signals' lifetime as well, increasing the number of needed registers.

A method of operating a dataprocessor device according to the invention is provided. The device comprises at least

-        a master controller for controlling operation of the device

-        a first functional unit, which includes a slave controller, the first functional unit being arranged for executing instructions of a first type corresponding to operations having a relatively long latency,

-        a second functional unit capable of executing instructions of a second type corresponding to operations having a relatively short latency. According to the method of the invention the first functional unit during execution of an instruction of the first type receives input data and provides output data, according to which method the output data is processed by the second functional unit during said execution and/or the input data is generated by the second functional unit during said execution.

The invention also provides for a method for compiling a program into a sequence of instructions for operating a processing device according to the invention. According to this method of compiling

-        a model is composed which is representative of the input/output operations involved in the execution of an instructions by a first functional unit,

-        on the basis of this model instructions for the one or more second functional units are scheduled for providing input data for the first functional unit when it is executing an instruction in which said input data is used and/or for retrieving output data from the first functional unit when it is executing an instruction in which said output data is computed.

These and other aspects of the invention are described in more detail with reference to the drawing. Therein

Figure 1 shows a data processing device,

Figure 2 shows an example of an operation which may be executed by the data processing device of Figure 1,

Figure 3A shows the signal flow graph (SFG) of the operation,

Figure 3B shows the operation's schedule and its time shape function,

Figure 4A schematically shows the operation of Figure 2,

Figure 4B shows a signal flow graph for schedulating execution of the operation of Figure 4A at a holdable custom functional unit (FU),

Figure 4C shows a signal flow graph for schedulating execution of the operation of Figure 4A at a custom functional unit (FU) which is not holdable,

Figure 5 shows a nested loop which includes the operation of Figure 2,

Figure 6A shows the traditional schedule of the nested loop of Figure 5 in a

5   SFG,

Figure 6B shows the schedule of said nested loop in a SFG according to the invention.

10   Figure 1 schematically shows a data processing device according to the invention. The data processing device at least comprises a master controller 1, a first functional unit 2 which includes a slave controller 20, a second functional unit 3. The two functional units 2, 3 share a memory 11 comprising a micro code as common memory means. The device is programmed for executing an instruction by the first functional unit 2, wherein

15   the execution of said instruction involves input/output operations by the first functional unit 2. The output data of the first functional unit 2 is processed by the second functional unit 3 during said execution and/or the input data is generated by the second functional unit 3 during said execution. In the embodiment shown the data processing device comprises further functional units 4, 5.

20   The embodiment of the data processing device shown in Figure 1 is characterized in that the first functional unit 2 is arranged for processing instructions of a first type corresponding to operations having a relatively large latency and in that the second functional unit 3 is arranged for processing instructions of a second type corresponding to operations having a relatively small latency.

25   As an example, the possible variation of FFT algorithms may be considered which can be implemented using an "FFT radix-4" FU. Then this custom FU can be re-used while the algorithm is modified from a decimation-in-time to a decimation-in-frequency FFT. The VLIW processor may perform other fine-grain operations while the embedded custom FU is busy with its coarse-grain operation. Therefore, the long latency coarse-grain operation

30   can be seen as a microthread [6] implemented on hardware, performing a separate thread while the remaining datapath's resources are performing other computations, belonging to the main thread.

Before introducing the scheduling problem, the Signal Flow Graph (SFG) [7][8][9] is defined as a way to represent the given application code. An SFG describes the

35   primitive operations performed in the code, and the dependencies between those operations.

Definition 1. Signal Flow Graph SFG.

A SFG is a 8-tuple $(V, I, O, T, E_d, E_s, w, \delta)$, where:

- V is a set of vertices (operations),
- I is the set of input,
5 - O is the set of output,
- $T \subseteq V \times I \cup O$ is the set of I/O operations' terminals,
- $E_d \subseteq T \times T$ is a set of data edges,
- $E_s \subseteq T \times T$ is a set of sequence edges, and
- $w : E_s \to Z$ is a function describing the timing delay

10 (in clock cycles) associated with each sequence edge.

- $\delta: V \to Z$ is a function describing the execution delay

(in clock cycles) associated with each SFG's operation.

In the definition of the SFG a distinction is made between directed data edges, and directed and *weighted* sequence edges. They impose different constraints in the scheduling problem

15 where "scheduling" is the task of determining for each operation $v \in V$, a start time s(v), subject to the precedence constraints specified by the SFG. Formally:

Definition 2. Traditional Scheduling Problem.

Given a SFG$(V, I, O, T, Ed, Es, w, \delta)$, find an integer labeling of the operations

$s: V \to Z^+$ where:

20

$$s(v_j) \geq s(v_i) + \delta(v_i) \qquad \forall i,j,h,k : ((v_i, o_h), (v_j, i_k)) \in E_d$$
$$s(v_j) \geq s(v_i) + w((t_i, t_j)) \qquad \forall i,j : (t_i, t_j) \in E_s$$

and the schedule's latency: $\max_{i=1..n}\{s(v_i)\}$ is minimum.

25

In the scheduling problem, as defined above, a single decision is taken for each operation, namely its start time. Because the I/O timeshape is not included in the analysis, no output signal is considered valid before the operation is completed. Likewise, the operation itself is started only if all input signals are available. This is surely a safe assumption, but allows no

30 synchronization between the operations' data consumption and production times and the start time of the other operations in the SFG.

Before formally stating the problem, an operation's timeshape is defined as follows:

Definition 3. Operation's timeshape

Given an SFG, for each operation $v \in V$, a *timeshape* is defined as the function

$\sigma : T_v \rightarrow Z^+$,

where:

5      $T_v = \{ t \in T \mid t=(v, p), \text{ with } p \in I \cup O \}$

is the set of I/O terminals for operation $v \in V$.

The number assigned to each I/O terminal models the delay of the I/O activity relatively to

the start time of the operation. Hence, for an operation of execution delay δ, the timeshape

10   function associates to each I/O terminal an integer value ranging from 0 to δ-1. An example

of operation's timeshape is depicted in Figure 3.

In the traditional scheduling problem, each operation is seen as atomic in the graph. In order

to exploit the notion of the operation's I/O timeshape, the scheduling problem is revisited.

Where a single decision was taken for each operation, now a number of decisions are taken.

15   Each scheduling decision is aimed to determine the start time of each I/O terminal belonging

to a given operation. Hence, the definition of the revisited scheduling problem taking into

account operations' timeshapes is the following:

Definition 4. I/O Timeshape Scheduling Problem:

20   Given a SFG and a timeshape functions for each operation $v \in V$ in the SFG, find an integer

labeling of the terminals $s : T \rightarrow Z^+$, where:

$s((v_j, i_k)) \geq s((v_i, o_h)) \qquad \forall i,j,h,k : (t(v_i, o_h), (v_j, i_k)) \in E_d$

$s(t_j) \geq s(t_i) + w((t_i, t_j)) \qquad \forall i,j : (t_i, t_j) \in E_s$

25

and the schedule's latency:

$\max_{i=1..n} \{ s(v_i) \}$ is minimum.

It is important to notice that, introducing the concept of timeshape, the operation's latency

30   function δ is not needed anymore and a scheduling decision is taken for each

operation's terminal. The schedule found must satisfy the constraints on data edges, sequence

edges, and respect the timing relations on the I/O terminals, as defined in the

timeshape functions. In order to exploit the I/O timeshape characteristic of operations, the timeshape function σ is translated in a number of sequence edges, added in the set E_s. These extra constraints impose that the start times of each I/O operation terminal, for any feasible schedule, are such that the timeshape of the original coarse-grain operations is respected.

5   The translation of the timeshape function into sequence edges is done in a different way depending on whether the FU implementing the coarse-grain operation, can or cannot be stopped during its computation. This will be discussed in more detail with reference to Figure 4. If the operation can be halted, then the timeshape of the operation can be stretched, provided that the concurrence and the sequence of the I/O terminals are kept. If

10   the unit cannot be halted then an extra constraint must be added in the graph, to make sure that not only the sequence but also the relative distance between I/O terminals is kept as imposed by timeshape function.

        By way of example two I/O terminals are considered which belong to the same original coarse-grain operation, namely t_1 and t_2. Then three different cases can happen:

15   1)   Concurrency

If two I/O terminals, t_1 and t_2, take place during the same cycle according to the timeshape of the coarse-grain operation, then two sequence edges are added. Those extra edges guarantee that the operations t_1 and t_2 in any feasible schedule, for the given SFG, will take place in the same cycle (e.g. in Figure 4B, o_1 and i_2).

20   If $\sigma(t_1) = \sigma(t_2)$ then (t_1,t_2), $(t_2,t_1) \in E_s$

with $w(t_1,t_2) = w(t_2,t_1) = 0$

According to the definition of the revisited scheduling problem, those two added edges impose that:

$s(t_1) \geq s(t_2)$ and $s(t_2) \geq s(t_1)$

25   he   Serialization (hold-able operation)

If two I/O terminals, t_1 and t_2, are not concurrent according to the coarse-grain operation's timeshape, then a sequence edge is added. This extra edge guarantees that the order of the two operations will be kept in any feasible schedule. Anyway, it allows that operation t_2 can be postponed relatively to operation t_1 (e.g. in Figure 4B, i_1 and i_2).

30   If $s(t_2) - s(t_1) = \lambda > 0$ then $(t_1,t_2) \in E_s$ with $w(t_1,t_2) = \lambda$

According to the definition of the revisited scheduling problem, this added edge imposes that:

$s(i_2) \geq s(i_1) + w(i_1,i_2) = s(i_1) + \lambda$

hence: $s(i_2) - s(i_1) \geq \lambda$

3)   Serialization (not hold-able operation)

The distance between the start times of the two I/O terminals, t_1 and t_2, is imposed, for any

feasible schedule, as defined by the coarse-grain timeshape (e.g. Figure 4C, i_1

and i_2). This is done adding two sequence edges: If $s(t_2) - s(t_1) = \lambda > 0$ then $(t_1, t_2), (t_2, t_1) \in E_s$

with $w(t_1, t_2) = \lambda$ and $w(t_2, t_1) = -\lambda$

According to the definition of the revisited scheduling problem, those two added edges

impose that:

$s(t_2) \geq s(t_1) + w(t_1, t_2) = s(t_1) + \lambda$

$s(t_1) \geq s(t_2) + w(t_2, t_1) = s(t_2) - \lambda$

From the last two equations, it follows that the difference in the starting time between t_1 and

t_2 is exactly equal to that imposed in the timeshape.

Hence:

$s(t_2) - s(t_1) = \lambda$

For each operation, the method adds a significant number of edges, in the order of $|I \cup O|^2$.

However, many of them can be pruned away, for instance introducing a partial order in the

set of the operation's terminals. The pruning step is mostly trivial and therefore, herewith not

described. Once the operations are described by their collection of I/O operations and the

sequence edges are added, the SFG is scheduled using known and traditional techniques.

Provided that the constraints due to the operations' timeshape are respected, the I/O terminals

of each operation are now de-coupled from each other and can be scheduled independently.

By way of example it is assumed that the given application is performing

intensively the "2Dtransform" function as shown in Figure 2. To make the example more

realistic, the function considered is performing a 2D graphic operation. It takes the vector

(x,y) and returns the vector (**X,Y**), according to the code as depicted in Figure 2. In order to

improve the processor's performance the "2Dtransform" is implemented in hardware on a

custom FU. Since the function is performed on hardware, it can be truly considered a single

coarse-grain operation. The signal flow graph for this function is depicted in Figure 3A. A

feasible internal schedule for the (coarse-grain) operation is depicted in Figure 3B, where one

adder and one multiplier, both with a latency of one cycle, are available within the custom

FU. The operation has four I/O terminals and it is performed by the custom FU in four clock

cycles, $\sigma = 0, \ldots, 3$.

In this example, although the FU is active during all the four cycles (Figure

3B), no I/O operation is performed in cycle 2. From the VLIW datapath, the internal

operations performed by the custom FU are not visible and only the I/O timeshape

is actually necessary to model the way the operation consumes and produces its data (Figure

5   3B).

The original coarse-grain operation in Figure 4A, whose content is now not

depicted, is re-modeled as a graph of four single cycle operations, each of them modeling an

I/O terminal. Sequence edges must be added to guarantee that the timeshape of the original

coarse-grain unit is respected in any possible feasible schedule. In the Figures the sequence

10   edges are indicated by dashed lines starting from a first operation and ending in an arrow at a

second operation. In Figure 4B, the derived SFG, modeling the behavior of a hold-able

custom FU, is shown. In particular, I/O terminals that were performed in different cycles,

according to the coarse-grain operation's timeshape, are serialized so that their order is

preserved. In said Figure for example an edge $w(i_1, i_2)$ having a value $\lambda = 1$ is present between

15   operations i_1 and i_2. Hence $s(i_2) \geq s(i_1) + w(i_1, i_2) = s(i_1) + \lambda$.. Concurrence of two or more I/O

terminals is kept as well. The time shape of Figure 4B for example comprises a first edge

$w(i_2, o_1)$ and a second edge $w(o_1, i_2)$ both having a value $\lambda = 0$ so that concurrence of the

operations i_2 and o_1 is garanteed. Hence, when a hold mechanism is available for the unit, the

scheduler can lengthen the coarse-grain operation moving I/O terminals apart from each

20   other, as far as the sequence edges are not violated. The effect on the hardware is that the FU

might be stalled to better synchronize data communicated to and from other operations.

Figure 4C shows the graph obtained by describing the coarse-grain operation

in I/O terminals when no hold mechanism is available for the custom FU. In this

case, the sequence edges added guarantee that the relative distance between any couple of I/O

25   terminals, in any feasible schedule, cannot be different from that imposed by the coarse-grain

operation's timeshape.

Now a code is considered where the function '2Dtransform' mapped on a

complex FU is used, as depicted in Figure 5. In this example, the "2Dtransform" operation is

part of a loop body, where other fine-grain operations, such as ALU operations and

30   multiplication's, are performed as well. It is supposed that the code is executed on a VLIW

processor containing in its datapath a multiplier, an adder and a "2Dtransform" FU.

The traditional schedule for the SFG of the above described loop body is

depicted in Figure 6A. The coarse-grain operation is regarded as "atomic" and no other

operation is executed in parallel with it. In Figure 6B the I/O schedule of the complex unit is

expanded and embedded in the loop body's SFG. The complex operation is executed concurrently with other fine-grain operations. According to the schedule, data is provided for the complex FU to the rest of the datapath and vice versa when actually needed, thereby reducing the schedule's latency. When some data is not available to the complex FU and the

5   computation cannot proceed further, the unit is halted (e.g. cycle 2 Figure 6B). The stall cycles are implicitly determined during the scheduling of the algorithm. Using the proposed solution, the latency of the algorithm is reduced from 10 to 8 cycles. The number of registers needed has decreased as well. The value produced in cycle 0 in Figure 6A has to be kept alive for two cycles, while the same signal in the schedule in Figure 6B is immediately used.

10   The proposed solution is efficient in terms of microcode area for the VLIW processor. The complex FU contains its own controller and the only task left to the VLIW controller is to synchronize the coarse-grain FU with the rest of the datapath resources. The only instructions that have to be sent to the unit are a start and a hold command. This can be encoded with few bits in the VLIW instruction word.

15   The VLIW processor can perform other operations while the embedded complex FU is busy with its computation.

The long latency unit can be seen as a micro-thread implemented on hardware, performing a task while the rest of the datapath is executing other computations using the rest of the datapath's resources.

20   The validity of the method has been tested using an FFT-radix4 algorithm as a case study. The FFT has been implemented for a VLIW architecture with distributed register files, synthesized using the architectural level synthesis tool "A|RT designer" from Frontier Design, running on a HP-UX machine. The radix-4 function, which constitutes the core of the considered FFT algorithm, processes 4 complex data values and 3 complex coefficients,

25   returning 4 complex output values. The custom unit "radix-4" contains internally an adder, a multiplier, and its own controller. The unit consumes 14 (real) input values and produces 8 (real) output values. Extra details of the "radix-4" FU are given in Table 1.

Table 1: The Radix4 Functional Unit.

30

	latency	internal registers	internal resources
Radix4 FU	26 cycles	16 (218 bits)	1 ALU, 1 MULT

Three different VLIW implementations are tested, as depicted in Table 2. The architectures "FFT_org" and "FFT_2ALU's") contain the same hardware resources but they differ in the coarseness of the operations that they can execute.

5   Table 2: The tested datapath architectures.

	Datapath Resources
FFT_org	1 ALU, 1MULT, 1 ACU, 1 RAM, 1 ROM
FFT_2ALU's	2 ALU, 1 MULT, 1 ACU, 1 RAM, 1 ROM
FFT_radix4	1 ALU, 1 ACU, 1 RADIX4, 1 RAM, 1 ROM

For each architecture instance, table 3 lists the performance of the implemented FFT radix4 algorithm in clock cycles and the dimension of the VLIW microcode memory, where the application's code is stored. If the first implementation ("FFT_org") is taken as a reference, it

10   can be observed in Table 3 that "FFT_2ALU's" presents the higher degree of parallelism and the best performance.

Table 3: Performance and microcode's dimension, experimental results.

	Performance (cycles)	Microcode (width x length)	Microcode width vs. original	Microcode n. bits
FFT_org	59701	76 * 82	100.0 %	6232
FFT_2ALU's	40145	95 * 61	125.0%	5795
FFT_radix4	49461	67 * 74	88.2%	4958

15

However, the extra ALU available in the datapath must be controlled directly by the VLIW controller, and a large increment in the microcode's instruction width is noticed. On the other side, "FFT_radix4" reaches performance which is in between the first two experiments, but a much narrower microcode memory is synthesized. Usually, the part of the code where the

20   parallelism is necessary is a small fraction of the entire code. If the FFT is a core functionality in a much longer application code then the microcode width, hence the ILP needed in "FFT_2ALU's", will not be exploited adequately in other portions of the code, leading to a waste of microcode area. "FFT_2ALU's" and "FFT_radix4" both offer 2 ALUs and a Multiplier in architecture for processing the critical FFT loop body, but fewer bits are

25   needed in the latter microcode to steer the available parallelism.

Table 4 lists, for each instance, the number of registers needed in the architecture. In particular, in the last architecture the total number of register is the sum of those present in the VLIW processor and those implemented within the "Radix4" unit. The experiments done confirm that scheduling the FFT SFG, exploiting the I/O timeshape of the

5   "Radix4" coarse-grain operation, reduces the number of needed registers.

Table 4: Register Pressure, experimental results.

	N. of registers	Registers total amount of bits
FFT_org	57	673
FFT_2ALU's	60	710
FFT_radix4	58 (42+16)	698 (481+218)

10   The method according to the invention allows for a flexible HW/SW partitioning where complex functions may be implemented in hardware as FUs in a VLIW datapath. The proposed "I/O timeshape scheduling" method allows for scheduling separately the start time of each I/O operation's event and, ultimately, to stretch the operation's timeshape itself to better adapt the operation with its surroundings. By using coarse-grain operations in VLIW

15   architectures, it is made possible to achieve high Instruction Level Parallelism without paying a heavy tribute in terms of microcode memory width. Keeping VLIW microcode width small is an essential requisite for embedded applications aiming at high performances and coping with long and complex program codes.

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